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Why do we need Pig?

- Writing native Map/Reduce is hard
 - Difficult to make abstractions
 - Extremely verbose
 - 400 lines of Java becomes < 30 lines of Pig
 - Joins are very difficult
 - A big motivator for Pig
 - Chaining together M/R jobs is tedious
 - Decouples logic from optimization
 - Optimize Pig and everyone benefits



Basically, everything about Java M/R is painful

Understanding a basic Pig script

```
This is a relation
                                 Loads a file into a relation, with a
(NOT a variable)
                                 defined schema
    students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int, dept:chararray);
                                            Filters out all rows that
                                            don't fulfill the predicate
  students filtered = FILTER students BY age >= 20;
     Loads another file
  dept info = LOAD 'dept info.txt' as (dept:chararray, credits:int);
                    Equi-joins two relations on
                    the dept column
  students dept = JOIN students BY dept, dept info BY dept;
                                                           Chooses what columns we care
 students proj = FOREACH students dept GENERATE
                                                           about and renames them
      students::first as first, students::last as last,
      students::dept as dept, dept info::credits as credits;
```

These are all relations! They represent rows of data, each row organized into columns

Basic Pig script (cont)

This makes a jar of User Defined Functions (UDFs) available for use register ClassesUdfs.jar;

Adds a column the the relation, the result of using a UDF

students_hours = FOREACH students_proj GENERATE *,

CalculateSemestersFromCredits(credits) as semesters;

Globally orders relation by the semesters column, in descending order

students_ordered = ORDER students_hours BY semesters DESC;

Stores the relation using a custom store function (could even store to a database)

STORE students_ordered INTO 'students_hours' USING CustomStudentsStoreFunc();



An even simpler example: Word count

 Can't have a tutorial on anything Hadoop related without word count

```
text = LOAD 'text' USING TextLoader(); Loads each line as one column tokens = FOREACH text GENERATE FLATTEN(TOKENIZE($0)) as word; wordcount = FOREACH (GROUP tokens BY word) GENERATE group as word, COUNT_STAR($1) as ct;
```



What Pig can do for you



What is Pig?

- Pig is a scripting language
 - No compiler
 - Rapid prototyping
 - Command line prompt (grunt shell)
- Pig is a domain specific language
 - No control flow (no if/then/else)
 - Specific to data flows
 - Not for writing ray tracers
 - For the distribution of a pre-existing ray tracer



What ISN'T Pig?

- A general framework for all distributed computation
 - Pig is MapReduce! Just easier
- A general purpose language
 - No scope
 - Minimal "variable" support
 - No control flow



What CAN Pig do?

- Ad hoc analysis
- Move and transform data (ETL)
 - Good integration with many systems
- Build regular reports
- Machine learning
- ...all on Terabytes and Petabytes of data!



What CAN'T Pig do?

- Iterative ML algorithms
 - Graph algorithms are generally slow
- Transcend Hadoop
 - Pig is still Hadoop
- Algorithms that converge
 - Without control flow can't check convergence
 - PigServer (Jython/Java) gives some control flow



What are my other options?

- Hive (SQL)
 - Pros
 - Leverages well-known syntax
 - JDBC means integration is easy
 - Cons
 - Complex transformations can be unwieldy
 - Less fine grained control



What are my other options? (cont)

- Cascading (Java API)
 - Various DSLs
 - Scalding (Scala API)
 - Cascalog (Clojure API)
 - Pros
 - All one language
 - The DSLs make Cascading simpler and more powerful
 - Easier to make converging algorithms
 - Cons
 - No optimization
 - A lower-level API



The Pig Object Model



Relations

- Fundamental building block
- Analogous to a table, not a variable

students = load 'students.txt' as (first:chararray, last:chararray, age:int, dept:chararray);

Defines a relation

A relation, as it appears in Excel

first	last	age	major
Willia	Cracknell	18	CS
Francesco	Corraro	21	English
Lino	Feddes	22	History
Wes	Knill	23	EE
Ellyn	Meyerhoefer	18	English
Warner	Caminita	24	Psych
Lucius	Orlosky	20	History
Del	Graefe	20	CS
Douglass	Adelizzi	23	CS
Lesley	Kellywood	20	Biology



Loading a file

- Save file as tab delimited text without headings
- Boot up Pig

```
$ bin/pig —x local 2012-08-15 14:29:04,968 [main] INFO org.apache.pig.Main - Apache Pig version 0.11.0-SNAPSHOT (r1373633) compiled Aug 15 2012, 14:20:59 2012-08-15 14:29:04,968 [main] INFO org.apache.pig.Main - Logging error messages to: /Users/jcoveney/workspace/berkeley_pig/pig_1345066144966.log 2012-08-15 14:29:04,983 [main] INFO org.apache.pig.impl.util.Utils - Default bootup file /Users/jcoveney/.pigbootup not found 2012-08-15 14:29:05,138 [main] INFO org.apache.pig.backend.hadoop.executionengine.HExecutionEngine - Connecting to hadoop file system at: file:/// grunt> students = LOAD 'students.txt' USING PigStorage('\t'); grunt> DUMP students;
```



A detour: firing up Pig

```
The version of Pig currently running
$ bin/pig -x local
2012-08-15 14:29:04,968 [main] INFO org.apache.pig.Main -
Apache Pig version 0.11.0-SNAPSHOT (r1373633) compiled Aug 15 2012, 14:20:59
2012-08-15 14:29:04,968 [main] INFO org.apache.pig.Main - Logging error messages to:
/Users/jcoveney/workspace/berkeley_pig/pig_1345066144966.log
2012-08-15 14:29:04,983 [main] INFO org.apache.pig.impl.util.Utils - Default bootup file /
Users/jcoveney/.pigbootup not found
2012-08-15 14:29:05,138 [main] INFO Errors will be logged here
org.apache.pig.backend.hadoop.executionengine.HExecutionEngine - Connecting to hadoop
file system at: file:///
Means we're working on local files
```

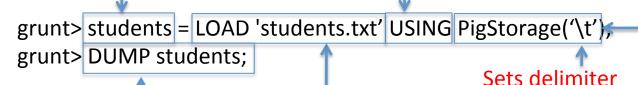
- Local mode is great for iterating quickly
 - Much faster than a cluster



Loading data

Declares the relation for the data

USING specifies the function that will load the data



which knows how to load rows from text

Refers to a Java function

DUMP keyword streams the relation to the console

Loads the data in this file

Output:

(Willia, Cracknell, 18, CS)

(Francesco, Corraro, 21, English)

(Lino, Feddes, 22, History)

(Wes, Knill, 23, EE)

(Ellyn, Meyerhoefer, 18, English)

(Warner, Caminita, 24, Psych)

(Lucius, Orlosky, 20, History)

(Del, Graefe, 20, CS)

(Douglass, Adelizzi, 23, CS)

(Lesley, Kellywood, 20, Biology)

It's the same as the excel file!

first	last	age	major
Willia	Cracknell	18	CS
Francesco	Corraro	21	English
Lino	Feddes	22	History
Wes	Knill	23	EE
Ellyn	Meyerhoefer	18	English
Warner	Caminita	24	Psych
Lucius	Orlosky	20	History
Del	Graefe	20	CS
Douglass	Adelizzi	23	CS
Lesley	Kellywood	20	Biology

Each row is a Tuple

Projection (aka FOREACH)

students = LOAD 'students.txt';

pruned = FOREACH students GENERATE \$0, \$2;

dump pruned;

Signifies the relation to transform

Here we define what each row in the new relation will be

we can refer to columns positionally

- Foreach means "do something on every row in a relation"
- Creates a new relation
- In this example, pruned is a new relation whose columns will be first name and age
- With schemas, can also use column aliases

Schemas and types

- Schema-less analysis is useful
 - Many data sources don't have clear types
- But schemas are also very useful
 - Ensuring correctness
 - Aiding optimization

```
students = LOAD 'students.txt'

as (first:chararray, last:chararray, age:int, dept:chararray);
```

The first half stays the same

- Schema gives an alias and type to the columns
 - Absent columns will be made null
- 7

Extra columns will be thrown out

We have the types, now use them!

This does the same as the previous example:

```
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int, dept:chararray); pruned = FOREACH students GENERATE first, age; DUMP pruned;
```

DESCRIBE prints the schema of a relation

```
grunt> describe pruned;

pruned: {first: chararray,last: chararray}

The relation being described The schema of the relation
```



Types and Schemas



Schemas vs. Types

- Schemas
 - A description of the types present
 - Used to help maintain correctness
 - Generally not enforced once script is run
- Types
 - Describes the data present in a column
 - Generally parallel Java types



Type Overview

- Pig has a nested object model
 - Nested types
 - Complex objects can contain other object
- Pig primitives mirror Java primitives
 - String, Int, Long, Float, Double
 - DataByteArray wraps a byte[]
 - Working to add native DateTime support, and more!



Complex types: Tuples

- Every row in a relation is a Tuple
- Allows random access
- Wraps ArrayList<Object>
- Must fit in memory
- Can have a Schema, but is not enforced
 - Potential for optimization



Complex types: Bags

- Pig's only spillable data structure
 - Full structure does not have to fit in memory
- Two key operations
 - Add a Tuple
 - Iterate over Tuples
- No random access!
- No order guarantees
- The object version of a relation
- Every row is a Tuple

Complex types: Maps

- Wraps HashMap<String, Object>
 - Keys must be Strings
- Must fit in memory
- Can be cumbersome to use
 - Value type poorly understood in script



More operators!



Filter

- A predicate is evaluated for each row
 - If false, the row is thrown out

```
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int, dept:chararray); under_20 = FILTER students BY age < 20;
```

```
(Willia, Cracknell, 18, CS) (Ellyn, Meyerhoefer, 18, English)
```

- Supports complicated predicates
 - Boolean logic
 - Regular expressions



See Pig documentation for more

Grouping

- Abstractly, GROUPing creates a relation with unique keys, and the associated rows
- Example: how many people in our data set are in each department?

key, so the schema of this Bag is the same as the

relation we grouped (ie students)

Second column is always named

after grouped relation, and is a Bag

Visualizing the group

```
(CS,{(Willia,Cracknell,18,CS),(Del,Graefe,20,CS),(Douglass,Adelizzi,23,CS)})
(EE,{(Wes,Knill,23,EE)})
(Psych,{(Warner,Caminita,24,Psych)})
(Biology,{(Lesley,Kellywood,20,Biology)})
(English,{(Francesco,Corraro,21,English),(Ellyn,Meyerhoefer,18,English)})
unique key → (History,{(Lino,Feddes,22,History),(Lucius,Orlosky,20,History)})

Every row
associated with
```

that unique key



An alternate visualization

Bags contain rows, so let's print it as such

```
students_grouped:
({(Willia,Cracknell,18,CS),
    (Del,Graefe,20,CS),
    (Douglass,Adelizzi,23,CS)})
({(Wes,Knill,23,EE)})
({(Warner,Caminita,24,Psych)})
({(Lesley,Kellywood,20,Biology)})
({(Francesco,Corraro,21,English),
    (Ellyn,Meyerhoefer,18,English)})
({(Lino,Feddes,22,History),
    (Lucius,Orlosky,20,History)})
```

students:

```
(Willia, Cracknell, 18, CS)
(Del, Graefe, 20, CS)
(Douglass, Adelizzi, 23, CS)
(Wes, Knill, 23, EE)
(Warner, Caminita, 24, Psych)
(Lesley, Kellywood, 20, Biology)
(Francesco, Corraro, 21, English)
(Ellyn, Meyerhoefer, 18, English)
(Lino, Feddes, 22, History)
(Lucius, Orlosky, 20, History)
```

They are the same!



Using GROUP: an example

 Goal: what % does each age group make up of the total?

```
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int, dept:chararray);
students grp = GROUP students BY age;
students ct = FOREACH students grp GENERATE group as age, COUNT STAR(students) as ct;
students total = FOREACH (GROUP students ct ALL) generate SUM(students ct.ct) as total;
students proj = FOREACH students join GENERATE
  students ct::age as age,
  (double)students ct::ct / (long)students total.total as pct;
(18,0.2)
(20,0.3)
(21,0.1) t It works! But how?
(22,0.1)
23.0.2)
```

Understanding GROUP

```
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int, dept:chararray);
                                                   Column "group" is
   students grp = GROUP students BY age;
                                                   the key we group on
   students_ct = FOREACH students_grp GENERATE group as age, COUNT_STAR(students) as ct;
                       Transformations can be nested
                                                               COUNT STAR gives us the
                                                               number of elements in the Bag
   students_total = FOREACH (GROUP students_ct ALL)
                                                               "students," which is the number
     GENERATE SUM(students_ct.ct) as total;
                                                               of rows with the same key
SUM UDF takes a Bag This syntax returns a
                                               ALL key means resulting relation will
and returns the sum
                                               be one row, the group column will be
                      Bag with just the
of the contents
                      specified column
                                               "all," and the second column will be
                                               every row in original relation
   students_proj = FOREACH students_join GENERATE
     students ct::age as age,
     (double)students_ct::ct / (long)students_total.total as pct;
```

cast for a decimal percent

Pig follows Java's math, so we This is a "scalar projection" of a relation. If a relation has one row with one value, we can cast it to that value.

Groups: a retrospective

- Grouping does not change the data
 - Reorganizes it based on the given key
 - Can group on multiple keys
- First column is always called group
 - A compound group key will be a Tuple ("group")whose elements are the keys
- Second column is a Bag
 - Name is the grouped relation
 - Contains every row associated with key



FLATTENing

- Flatten is the opposite of group
- Turns Tuples into columns
- Turns Bags into rows



Flattening Tuples

```
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int, dept:chararray);
students_grouped = GROUP students BY (age, dept); We are grouping on multiple keys
students ct = FOREACH students_grouped GENERATE group, COUNT(students) as ct;
DESCRIBE students ct;
                                                       Irrespective of the number of group
students_ct: {group: (age: int,dept: chararray),ct:long} keys, there is always one first
                                                       column, and it is called "group"
dump students ct;
                                                   This is the schema of a Tuple.
((18,CS),1)
             ___ This is how a Tuple looks.
                                                   students ct has TWO columns,
((18, English), 1) Flatten let's us un-nest the
                                                    one of which is a Tuple
((20,CS),1)
              columns it contains
((20,Biology),1)
((20, History), 1)
((21,English),1)
((22, History), 1)
((23,CS),1)
(23,EE),1)
```

(24,Psych),1)

Flattening Tuples (cont)

```
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int, dept:chararray);
students grouped = GROUP students BY (age, dept);
students ct = FOREACH students_grouped GENERATE FLATTEN(group) as (age, dept),
                                                 COUNT(students) as ct;
DESCRIBE students ct;
                                              Flatten un-nests the The same re-aliasing
students_ct: {age: int,dept: chararray, ct:long}
                                              compound key
                                                                   statement as before
                                                                   works when working
dump students_ct;
                                  The columns have been
                                                                   with Tuples
(18,CS,1) The values are the
                                  brought down a level.
(18,English,1) same, but now
                                  students ct now has
(20,CS,1)
          there are two
                                  three columns
(20,Biology,1) columns instead of
(20, History, 1) a Tuple
(21,English,1)
(22, History, 1)
(23,CS,1)
23,EE,1)
```

24,Psych,1)

Flattening Bags

- Syntax is the same as Flattening Tuples, but the idea is different
- Tuples contain columns, thus flattening a
 Tuple turns one column into many columns
- Bags contain rows, so flattening a Bag turns one row into many rows



Returning to an earlier example

```
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int,
dept:chararray);
students grouped = GROUP students BY dept;
students proj = FOREACH students_grouped GENERATE students;
students grouped:
({(Willia, Cracknell, 18, CS),
                                             (Willia, Cracknell, 18, CS)
 (Del, Graefe, 20, CS),
                                             (Del, Graefe, 20, CS)
 (Douglass, Adelizzi, 23, CS)))
                                             (Douglass, Adelizzi, 23, CS)
({(Wes,Knill,23,EE)})
                                             (Wes, Knill, 23, EE)
({(Warner, Caminita, 24, Psych)})
                                             (Warner, Caminita, 24, Psych)
({(Lesley,Kellywood,20,Biology)})
                                             (Lesley, Kellywood, 20, Biology)
({(Francesco, Corraro, 21, English),
                                             (Francesco, Corraro, 21, English)
 (Ellyn, Meyerhoefer, 18, English)))
                                             (Ellyn, Meyerhoefer, 18, English)
({(Lino,Feddes,22,History),
                                             (Lino, Feddes, 22, History)
 (Lucius, Orlosky, 20, History)})
                                             (Lucius, Orlosky, 20, History)
```

Data is the same, just with different nesting. On the left, the rows are divided into different Bags.



Flattening Bags (cont)

The schema indicates what is going on

```
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int,
   dept:chararray);
                                                                       Notice that the two
   DESCRIBE students;
                                                                       schemas are basically the
Posttudents: her stight a ray lasting age: int, dept: chararray
                                                                       same, except in the
   students grouped = GROUP students BY dept;
                                                                       second case the rows are
   students proj = FOREACH students grouped GENERATE students;
                                                                       contained inside of a Bag
   DESCRIBE students_proj;
   students proj: {students: {(first: chararray,last: chararray,age: int,dept: chararray)}}
   students flatten = FOREACH students proj GENERATE FLATTEN(students);
   DESCRIBE students_flatten;
 students_flatten: {students::first: chararray,students::last: chararray,students::age:
   int,students::dept: chararray}
```

Group goes from a flat structure to a nested one



So let's flatten the Bag already

 Now that we've seen grouping, there's a useful operation

```
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int,
dept:chararray);
students grouped = GROUP students BY dept;
students_proj = FOREACH students_grouped GENERATE
DUMP students projLATTEN(students) as (first, last, age, dept);
(Willia, Cracknell, 18, CS)
                                       Same re-aliasing statement as before
(Del, Graefe, 20, CS)
                                       works with the rows resulting from
(Douglass, Adelizzi, 23, CS)
                                       flattened Bags
(Wes,Knill,23,EE)
(Warner, Caminita, 24, Psych)
(Lesley, Kellywood, 20, Biology)
(Francesco, Corraro, 21, English)
                                       — It worked! We have the original rows
(Ellyn, Meyerhoefer, 18, English)
Lino, Feddes, 22, History)
Lucius, Orlosky, 20, History)
```

Fun with flattens

The GROUP example can be done with flattens

```
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int,
dept:chararray);
students grp = GROUP students BY age;
students ct = FOREACH students grp GENERATE
 group as age,
 COUNT STAR(students) as ct;
students total = FOREACH (GROUP student [attAnt) of the African rows
 SUM(students ct.ct) as total;
students_proj = FOREACH students_total GENERATE scalar projection
  age,
  (double) ct / total as pct;
```

Joins

- A big motivator for Pig easier joins
- Compares relations using a given key
- Output all combinations of rows with equal keys
- See appendix for more variations

```
dept_info = LOAD 'dept_info.txt' as (dept:chararray, credits:int);
DUMP dept_info;
```

(CS, 1000)

(EE,1000)

(Psych, 10)

(Biology, 20)

(English,1)

History,3)

Let's introduce a new data source related to our previous one. This one connects a department to its required credits

Joins (cont)

How many credits does each student need?

```
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int, dept:chararray);
dept_info = LOAD 'dept_info.txt' as (dept:chararray, credits:int);
students_dept = JOIN students_BY dept, dept_info_BY dept;
This signifies the join key
(the keys whose equality
will be tested)
students_dept: {students::first: chararray,students::last: chararray,students::age:
int,students::dept: chararray,dept_info::credits: int}
```

- Joined schema is concatenation of joined relations' schemas
- Relation name appended to aliases in case of ambiguity
 - In this case, there are two "dept" aliases

Order by

- Order by globally sorts a relation on a key (or set of keys)
- Global sort not guaranteed to be preserved through other transformations
- A store after a global sort will result in one or more globally sorted part files

Extending Pig: UDFs



Extending Pig

 UDF's, coupled with Pig's object model, allow for extensive transformation and analysis

```
UDF must be compiled and packaged into a jar, then registered.
```

```
register JarContainingUdf.jar
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int, dept:chararray);
students_names = FOREACH students GENERATE myudfs.ToUpperCase(first);
```

UDF is invoked by name, including package (case-sensitive!)



What's a UDF?

- A User Defined Function (UDF) is a java function implementing EvalFunc<T>, and can be used in a Pig script
 - Additional support for functions in Jython, JRuby,
 Groovy, and Rhino (experimental)
- Much of the core Pig functionality is actually implemented in UDFs
 - COUNT in the previous example
 - Useful for learning how to implement your own
 - src/org/apache/pig/builtin has many examples



Types of UDFs

- EvalFunc<T>
 - Simple, one to one functions
- Accumulator<T>
 - Many to one
 - Left associative, NOT commutative
- Algebraic<T>
 - Many to one
 - Associative, commutative
 - Makes use of combiners
- All UDFs must returns Pig types
 - Even intermediate stages

EvalFunc<T>

- Simplest kind of UDF
- Only need to implement an "exec" function
- Not ideal for "many to one" functions that vastly reduce amount of data (such as SUM or COUNT)
 - In these cases, Algebraics are superior
- src/org/apache/pig/builtin/TOKENIZE.java is a nontrivial example



A basic UDF

```
package myudfs;
import org.apache.EvalFunc;
public class ToUpperCase extends EvalFunc String > {
    public String exec(Tuple input); {
        String inp = (String) input.get(0);
        return inp.toUpperCase();
    }

Input is always a Tuple (ie a row). Thus,
    UDF input is also untyped
    return inp.toUpperCase();
}

Input Tuple is untyped, so we must cast
    entries ourselves.
```



What happens when you run a script?



The DAG

- Pig script results in 1+ MapReduce jobs
- Graph of dependencies between these jobs is a directed acyclic graph (DAG)
 - http://www.github.com/twitter/ambrose is a great tool for visualizing a Pig script's DAG
- DAG is a result of Map/Reduce barriers



What is a Map/Reduce barrier

- A Map/Reduce barrier is a part of a script that forces a reduce stage
 - Some scripts can be done with just mappers

```
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int, dept:chararray);
students_filtered = FILTER students BY age >= 20;
students_proj = FOREACH.students_filtered GENERATE last, dept;
— But most will need the full MapReduce cycle
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int, dept:chararray);
students_grouped = GROUP students BY dept;
students_proj = FOREACH students_grouped GENERATE group, COUNT(students);
— The group is the difference, a "map reduce"
```

barrier" which requires a reduce step

M/R implications of operators

- What will cause a map/reduce job?
 - GROUP and COGROUP
 - JOIN
 - Excluding replicated join (see Appendix)
 - CROSS
 - To be avoided unless you are absolutely certain
 - Potential for huge explosion in data
 - ORDER
 - DISTINCT
- What will cause multiple map reduce jobs?
 - Multiple uses of the above operations
 - Forking code paths



Making it more concrete

First step is identifying the M/R barriers

```
Job#
     students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int, dept:chararray);
     students grp = GROUP students BY dept;
     students ct = FOREACH students grp GENERATE group as dept, COUNT(students) as ct;
     students_grp2 = GROUP students_ct BY ct;
     students_ct2 = FOREACH students_grp2 GENERATE group, COUNT(students_ct);
     STORE students ct2 INTO 'histogram';
     dept ct = FILTER student_ct BY ct > 2;
     students big dept = JOIN students BY dept, dept ct BY dept;
     students_big_dept_proj = FOREACH students_big_dept GENERATE
       students::first as first, students::last as last, students::age as age, students::dept as dept;
     STORE students big dept proj INTO 'students in big departments';
```



A DAG example

Job 1

Map LOAD 'students.txt'

GROUP students BY dept

Reduce GENERATE group as dept, COUNT(students) as ct

Map FILTER student_ct BY ct > 2

JOIN students BY dept, dept_ct BY dept

Reduce GENERATE first, last, age, dept;

STORE students big dept proj

Job 2

GROUP students_ct BY ct

GENERATE group, COUNT(students_ct)

STORE students ct2

Job 3



Life lessons



How to make good Pig scripts

- Project early, project often
 - Always better to be more explicit
 - Reduces the amount of information being shuffled
- Explicitly name things
 - Pig is smart about it's schemas, but making things explicit will make scripts much more readable



How to make good Pig scripts (cont)

- Don't reuse relation names
 - Makes troubleshooting failing M/R jobs harder
- For UDFs, implement Accumulator and Algebraic if possible
 - Unless it is just a simple one to one EvalFunc
- General life lesson: nobody likes unreadable, terse code, no matter how powerful the language. Be explicit!



Getting help from the Pig community

- Read the docs! http://pig.apache.org has more formal information on all of this
- There are a bunch of Pig experts that want to help you: <u>user@pig.apache.org</u>
- When having issues, please include the script
 - Ideally, a smaller script that isolates your error
 - Example data even more ideal
- If you found a bug, file a bug report!
 - https://issues.apache.org/jira/browse/PIG



Appendix



More on projections

- Projection reduces the amount of data being processed
 - Especially important between map and reduce stages when data goes over the network

```
rel1 = LOAD 'data1.txt' as (a, b, c, ... lots of columns ..);
rel2 = LOAD 'data2.txt' as (a, ... lots of columns ..);
rel3 = JOIN rel1 BY a, rel2 BY a;
rel4 = FOREACH rel3 GENERATE rel1::a, rel1::b, rel1::c;
```

In this case, all of the columns in relation1 and relation2 would be so across the network! Pig tries to optimize this, but sometimes fails. Remember: be explicit!

```
rel1 = LOAD 'data1.txt' as (a, b, c, ... lots of columns ..);
rel1_proj = FOREACH rel1 GENERATE a, b, c;
rel2 = LOAD 'data2.txt' as (a, ... lots of columns ..);
be shuffled, a code more extend = JOIN rel1_proj BY a, rel2_proj BY a;
rel4 = FOREACH rel3 GENERATE rel1_proj::a, rel1_proj::b, rel1_proj::c;
```

This ensures extra data won't be shuffled, and makes your code more explicit

Scalar projection

- All interaction and transformation is in Pig is done on relations
- Sometimes, we want access to an aggregate
 - Scalar projection allow us to use intermediate aggregate results in a script

```
students = LOAD 'students.txt';
count = FOREACH (GROUP students ALL) GENERATE COUNT_STAR($1) as ct;
```

(Incidentally, this is the pattern for counting a relation)

This will make the column ct in count available as a long

proj = foreach students generate *, (long) count.ct;



If the relation has more than one row or the specified column isn't of the right type, it will error out

More on SUM, COUNT, COUNT_STAR

 In general, SUM, COUNT, and other aggregates implicitly work on the first column

```
rel1 = LOAD 'data1.txt' as (x:int, y:int, z:int);
rel2 = GROUP rel1 ALL;
rel3 = FOREACH rel2 GENERATE SUM(rel1);
rel4 = FOREACH rel2 GENERATE SUM(rel1.x);
rel5 = FOREACH rel2 GENERATE SUM(rel1.y);
These won't
```

- COUNT counts only non-null fields
- COUNT_STAR counts all fields

```
rel6 = FOREACH rel2 GENERATE COUNT(rel1); This counts only non-null values of x
```

Whereas this counts the number of values in the relation

More on sorting

- Sorting is a global operation, but can be distributed
- Must approximate distribution of the sort key
- Imagine evenly distributed data between 1 and 100. With 10 reducers, can send 1-10 to computer 1, 11-20 to computer 2, and so on. In this way, the computation is distributed but the sort is global
- Pig inserts a sorting job before an order by to
 estimate the key distribution

More on spilling

- Spilling means that, at any time, a data structure can be asked to write itself to disk
- In Pig, there is a memory usage threshold
- This is why you can only add to Bags, or iterate on them
 - Adding could force a spill to disk
 - Iterating can mean having to go to disk for the contents



Flattening multiple Bags

- The result from multiple flatten statements will be crossed
 - To only select a few columns in a Bag, syntax is bag_alias.(col1, col2)

```
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int, dept:chararray);
students_grouped = GROUP students BY dept;
students_grouped = FOR EACH students.grouped GENERATE FLATTEN(students.first),
    running them to understand what's going on FLATTEN(students.age);
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int, dept:chararray);
students_grouped = GROUP students BY dept;
students_proj = FOREACH students_grouped GENERATE FLATTEN(students.(first, age));
    This is the proper way to choose a specific set of columns from a bag
```

Nested foreach

- An advanced, but extremely powerful use of FOREACH let's a script do more analysis on the reducers
- Imagine we wanted the distinct number of ages per department

```
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int, dept:chararray);
students_grouped = GROUP students BY dept;
unique_ages = FOREACH students_grouped {
   dst = DISTINCT students.age;
   GENERATE group as dept, FLATTEN(dst) as age;
}
this creates a Bag, dst, which is the istinct of the Bag students.age

| Students_grouped {
   Within the nested block can use a subset of Pig commands (FILTER, DISTINCT, ORDER BY) to manipulate the bag of rows associated with the group key
```

Nested foreach: be careful!

 Is very useful, but since computation is done in memory, can blow up

```
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int,
dept:chararray);
students grouped = GROUP students BY dept;
unique ages = FOREACH students grouped {
  dst = DISTINCT students.age;
  GENERATE group as dept, FLATTEN(dst) as age; Is the same as...
students = LOAD 'students.txt' as (first:chararray, last:chararray, age:int,
dept:chararray);
students dst = FOREACH (GROUP students BY (dept, age))
  GENERATE FLATTEN(group) as (dept, age);
unique ages = FOREACH (GROUP students dst BY dept)
  GENERATE group as dept, FLATTEN($1.age) as age;
 slower, requiring 2 MR jobs
```

Join optimizations

- Pig has three join optimizations. Using them can potentially make jobs run MUCH faster
- Replicated join
 - a = join rel1 by x, rel2 by x using 'replicated';
- Skewed join
 - a = join rel1 by x, rel2 by x using 'skewed';
- Merge join
 - a = join rel1 by x, rel2 by x using 'merge';



Replicated join

- Can be used when:
 - Every relation besides the left-most relation can fit in memory
- Will invoke a map-side join
 - Will load all other relations into memory in the mapper and do the join in place
- Where applicable, massive resources savings



Skewed join

- Useful when one of the relations being joined has a key which dominates
 - Web logs, for example, often have a logged out user id which can be a large % of the keys
- The algorithm first samples the key distribution, and then replicate the most popular keys
 - Some overhead, but worth it in cases of bad skew
- Only works if there is skew in one relation
 - If both relations have skew, the join degenerates to a cross, which is unavoidable

Merge join

- This is useful when you have two relations that are already ordered
 - Cutting edge let's you put an "order by" before the merge join
- Will index the blocks that correspond to the relations, then will do a traditional merge algorithm
- Huge saving when applicable



Tying it to Hadoop



So what's actually going on?

- So now we can write a script, but we haven't really talked about how Pig executes a script "under the hood"
- Pig puts together a Map/Reduce job based on the script that you give it



What happens when run a script?

- First, Pig parses your script using ANTLR
- The parser creates an intermediate representation (AST)
- The AST is converted to a Logical Plan
- The Logical Plan is optimized, then convert to a Physical Plan
- The Physical Plan is optimized, then converted to a series of Map/Reduce jobs



Wait...what?

- Layers of abstractions are very useful for performing optimizations at various levels
- Logical Plan
 - High level description of the computation
- Physical Plan
 - Pipeline that performs the computation
- Map/Reduce Job
 - Graph of jobs that actually run on the cluster



Logical Plan

- High level description of the data flow
- Describes the computation that will be done,
 without implementation
- Can be optimized
 - Column pruning
 - Throws out any unused columns
 - Filter push-down
 - Push filters as high as possible in order to reduce data being processed and shuffled



Physical plan

- Physical description of the computation
- Creates a usable pipeline
 - Pipeline usable independent of M/R
 - Could use this pipeline and target other processing frameworks
- This stage can also be optimized
 - In memory aggregation instead of combiners



MapReduce plan

- The logical and physical plans are ideally divorced from the details of running on Hadoop
 - This is not always the case, but it's close
- MRCompiler breaks the physical plan into a DAG of M/R jobs
- Can be optimized as well
 - Combining multiple M/R jobs into one



How do I see the plans?

- Pigs "explain <relation>" will print the three plans that Pig generates
 - Extremely useful for debugging
 - Can be a bit advanced, beyond scope of this presentation
 - Pig listserv can be very helpful



Advanced UDFs



Accumulator<T>

- Used when the input is a large bag, but order matters
 - Allows you to work on the bag incrementally, can be much more memory efficient
- Difference between Algebraic UDFs is generally that you need to work on the data in a given order
 - Used for session analysis when you need to analyze events in the order they actually occurred
- src/org/apache/pig/builtin/COUNT.java is an example
 - Also implements Algebraic (most Algebraic functions are also Accumulative)



Algebraic

- Commutative, algebraic functions
 - You can apply the function to any subset of the data (even partial results) in any order
- The most efficient
 - Takes advantage of Hadoop combiners
 - Also the most complicated ☺
- src/org/apache/pig/builtin/COUNT.java is an example



Algebraic: initial function

- Accepts a Tuple with a Bag with one row
 - Complicated, but that's how it is
- Makes no assumptions on order of execution
 - Instantiation may be executed on 0 or more rows
- Must return a Tuple containing valid Pig data type
 - EvalFunc<Tuple>



Algebraic: intermed function

- Accepts Tuple which contains a Bag of results from the initial or intermediate function
 - Must be able to accept it's own output
- Makes no assumptions on how many elements that Bag will contain
- Might not be called at all
- Must return a Tuple containing valid Pig data types



Algebraic: final function

- Responsible for returning the final aggregate
- Accepts Tuple which contains a Bag of results from initial or intermediate function
- Only invoked once
- Must return a valid Pig data type
 - EvalFunc<T>



Algebraic: Sum example

This abstract class lets us just implement Algebraic

```
public SimplerSum extends AlgebraicEvalFunc<Long> {
                                                          This shared class makes Tuples
  private static final TupleFactory mTupleFactory = TupleFactory.getInstance();
  public String getInitial() { return Initial.class.getName(); }
                                                                  These functions return
  public String getIntermed() { return Initial.class.getName(); }
                                                                  the classes that will do
  public String getFinal() { return Final.class.getName(); }
                                                                  the processing
                           In this case, the initial and intermediate
                          functions are the same
  private static long sumBag(DataBag inp) {
    long sum = 0;
                                                                 This helper function actually
    for (Tuple t : inp) { sum += ((Number)t.get(0)).longValue(); } does the summing
    return sum;
  // implementation continued on next slide
```

Algebraics: Sum example (cont)

```
public SimplerSum implements Algebraic {
  // assumes previous slide
                                              Both classes extend the EvalFunc class
  static class Initial extends EvalFunc<Tuple> {
    public Tuple exec(Tuple input) throws IOException {
       Tuple out = mTupleFactory.newTuple(1);
                                                   In this case, the Initial function can also
      out.set(0, sumBag((DataBag)input.get(0)));
                                                   work on collections of it's own output
      return out;
  static class Final extends EvalFunc<Long> {
    public Long exec(Tuple input) throws IOException {
      return sumBag((DataBag)input.get(0));
                                                 Only difference is that it returns the
                                                 actual Long instead of wrapping it
```

Other extension interfaces

- Pig has other interfaces used in advanced cases
- AlgebraicEvalFunc<K>
 - Implement Algebraic, get Accumulator for free
- AccumulatorEvalFunc<K>
 - Implement Accumulator, get EvalFunc free
- TerminatingAccumulator<K>
 - Accumulator which can terminate early
 - Can save a significant amount of processing
- IteratingAccumulatorEvalFunc<K>
 - Provides a more human accumulator interface (at the cost of some potential overhead)

See the documentation!

- The documentation has information on all of these topics and more, including
 - Outer joins
 - Cogroups
 - Macros
 - Parameters
 - More in-depth Accumulator and Algebraic
- Also be sure to look at the source of the builtin functions (as well as those in Piggybank)
 - This is a huge benefit of open source!